

CS 5244: Introduction to Cyber Physical Systems

Unit 5: Memory Architectures (Ch. 8)

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**Acknowledgement: The instructor thanks Profs. Edward A. Lee & Sanjit
A. Seshia at UC Berkeley for sharing their course materials**

Memory Architecture: Issues

- Types of memory
- Stack
- Caches
- Scratchpad memories
- Absolute and relative addresses
- Virtual memory
- Heaps
 - allocation/deallocation
 - fragmentation
 - garbage collection
- Segmented memory spaces
- ...

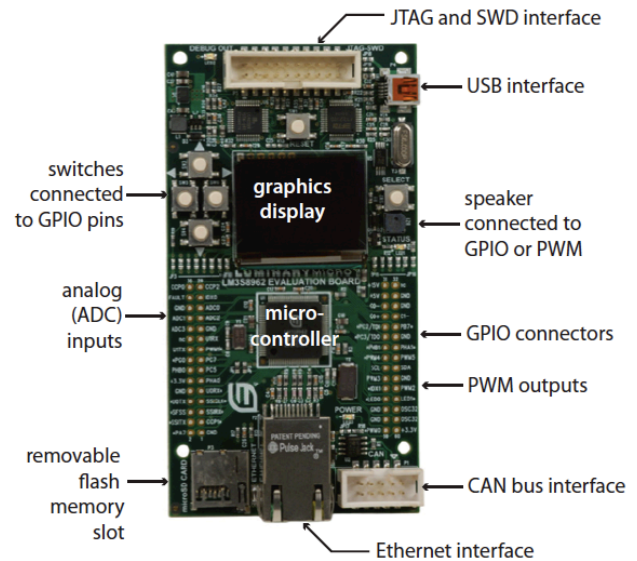
These issues loom larger in embedded systems than in general-purpose computing.

Specific Examples

To be concrete:

- A low-end example:
8-bit microcomputer

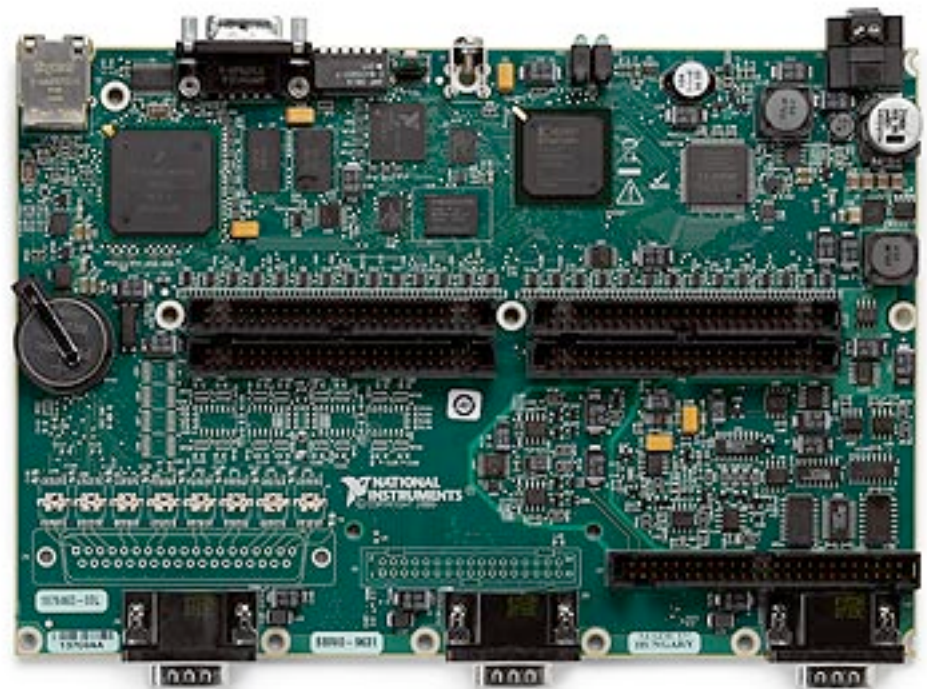
- A medium-end example:
32-bit microcomputer



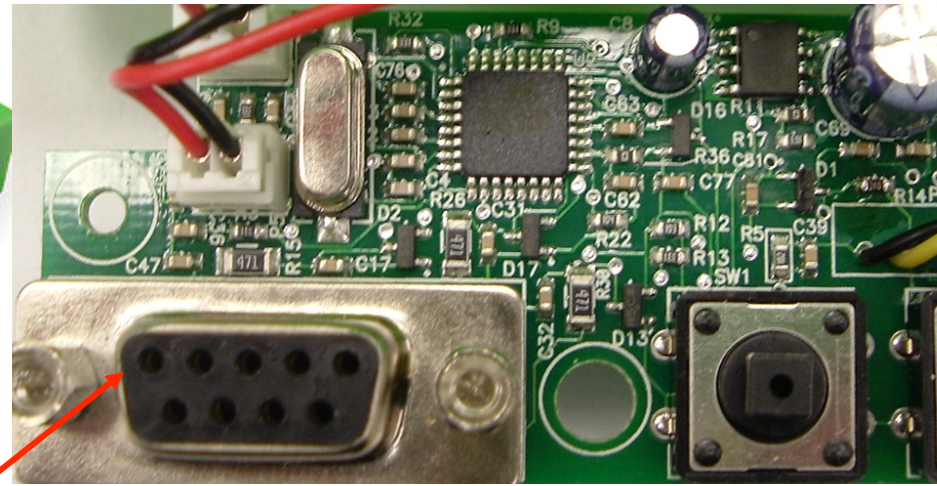
Specific Examples (Continued)

Single-Board RIO (National Instruments)

- Xilinx FPGA
 - In our lab: preconfigured with a 32-bit MicroBlaze microprocessor running without an operating system (“bare iron”).
- PowerPC processor (Freescale MPC5200)
 - In our lab: running VxWorks RTOS (real-time operating system) or LabVIEW Embedded.

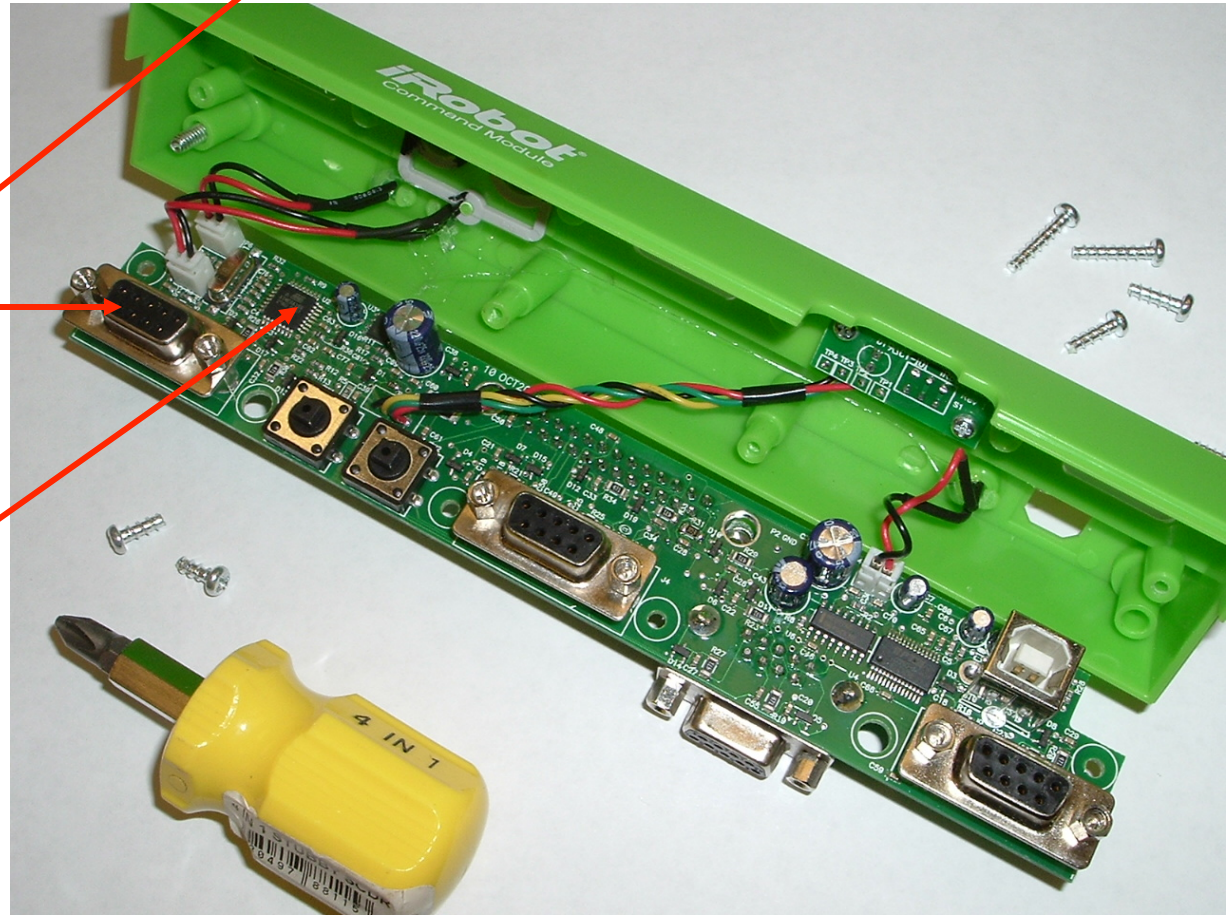


The *iRobot Create* Command Module



9-pin I/O port

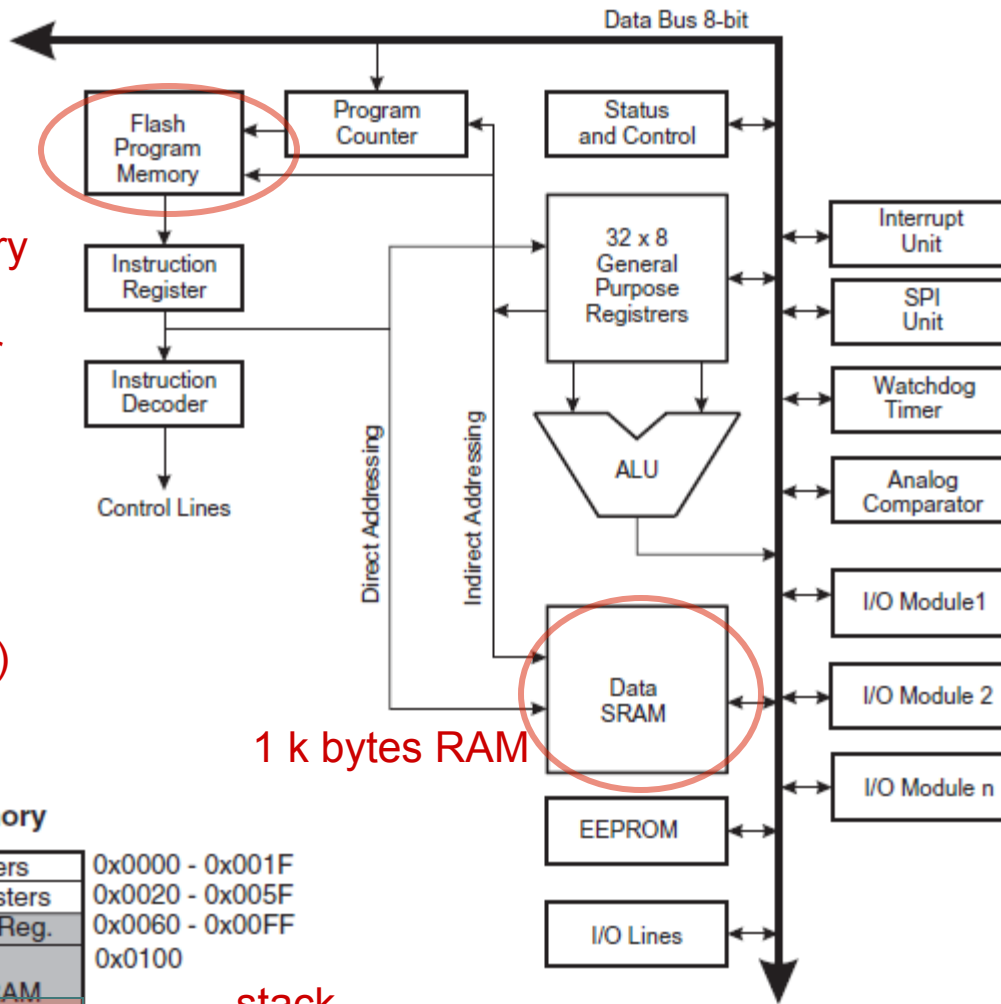
Atmel ATmega 168
Microcontroller



ATMega168 Memory Architecture

An 8-bit microcontroller with 16-bit addresses

iRobot command module has 16K bytes flash memory (14,336 available for the user program. Includes interrupt vectors and boot loader.)



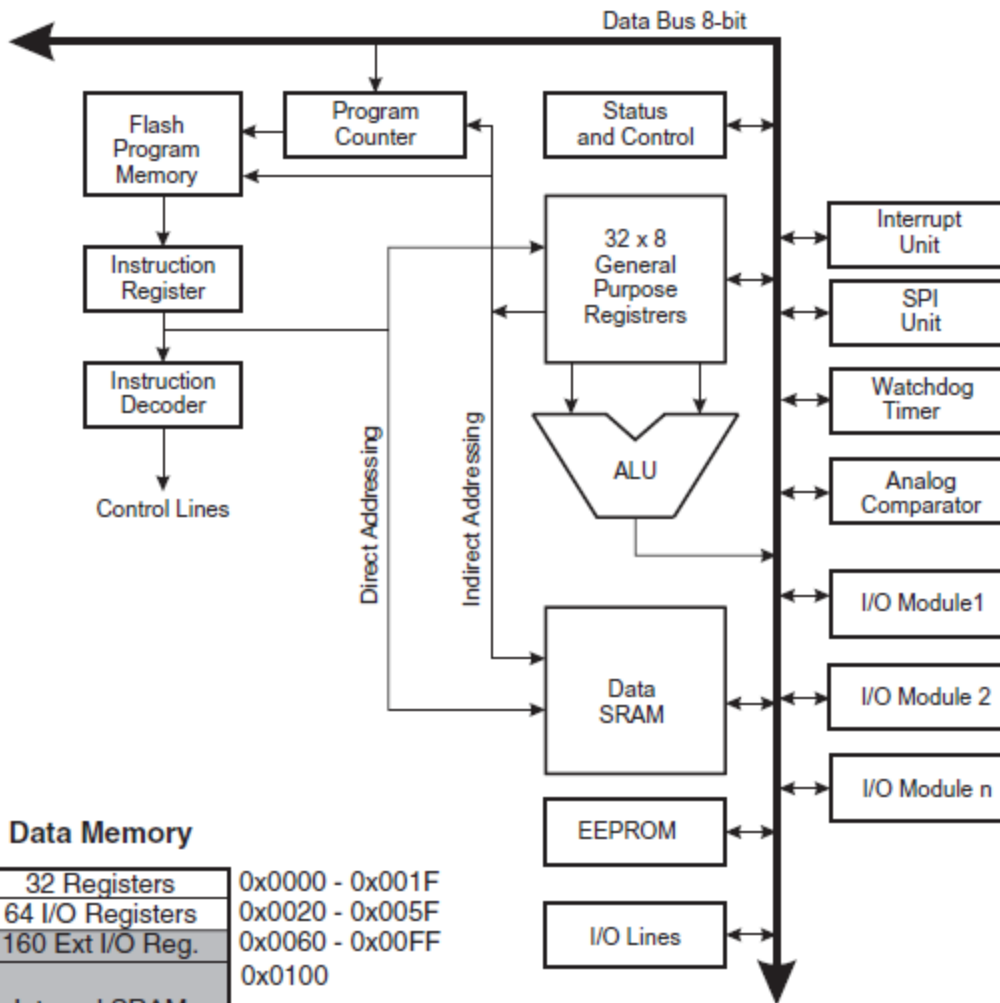
1 k bytes RAM

Data Memory	
32 Registers	0x0000 - 0x001F
64 I/O Registers	0x0020 - 0x005F
160 Ext I/O Reg.	0x0060 - 0x00FF
Internal SRAM	0x0100
(512/1024/1024 x 8)	
stack	0x02FF/0x04FF/0x04FF

Example of a microcontroller architecture. Used in iRobot command module.

- Additional I/O on the command module:
- Two 8-bit timer/counters
 - One 16-bit timer/counter
 - 6 PWM channels
 - 8-channel, 10-bit ADC
 - One serial UART
 - 2-wire serial interface

Understanding the memory architecture



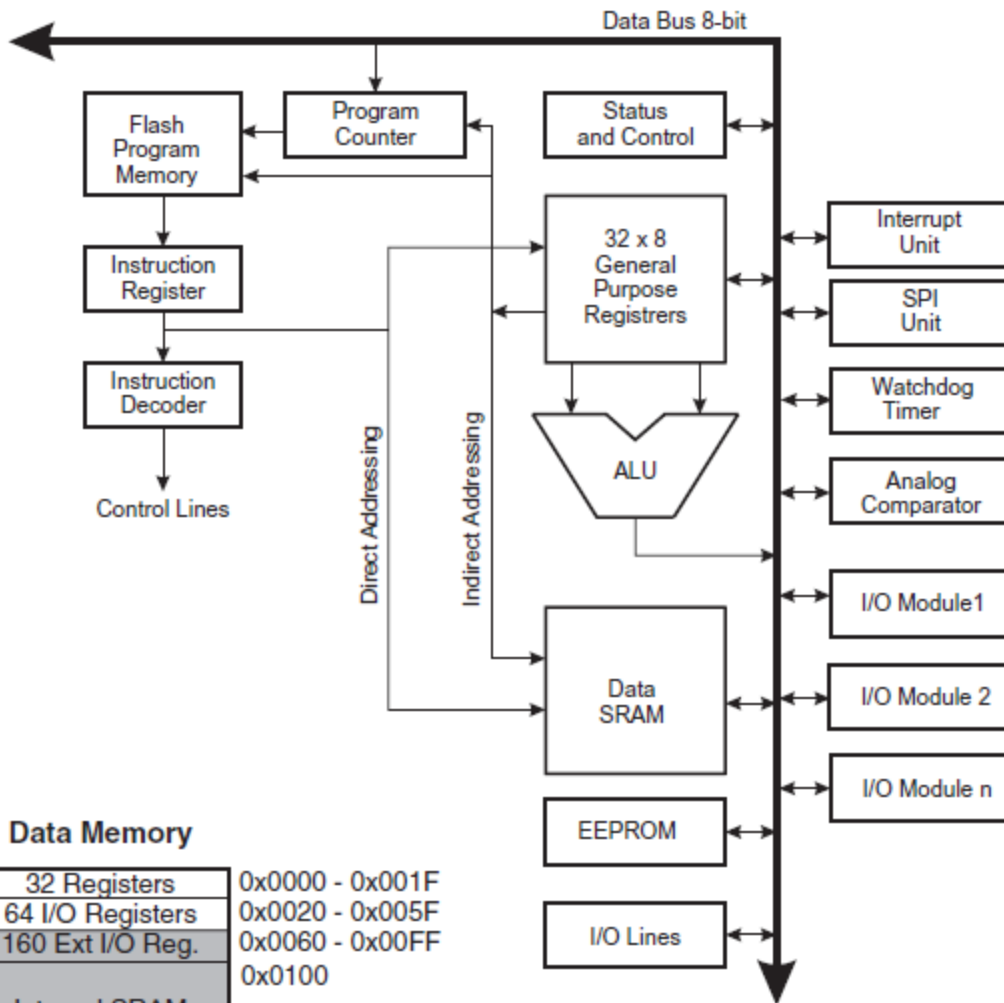
Data Memory

32 Registers	0x0000 - 0x001F
64 I/O Registers	0x0020 - 0x005F
160 Ext I/O Reg.	0x0060 - 0x00FF
	0x0100
Internal SRAM (512/1024/1024 x 8)	0x02FF/0x04FF/0x04FF

What is meant by the following C code:

```
char x;
x = 0x20;
```

Understanding the memory architecture



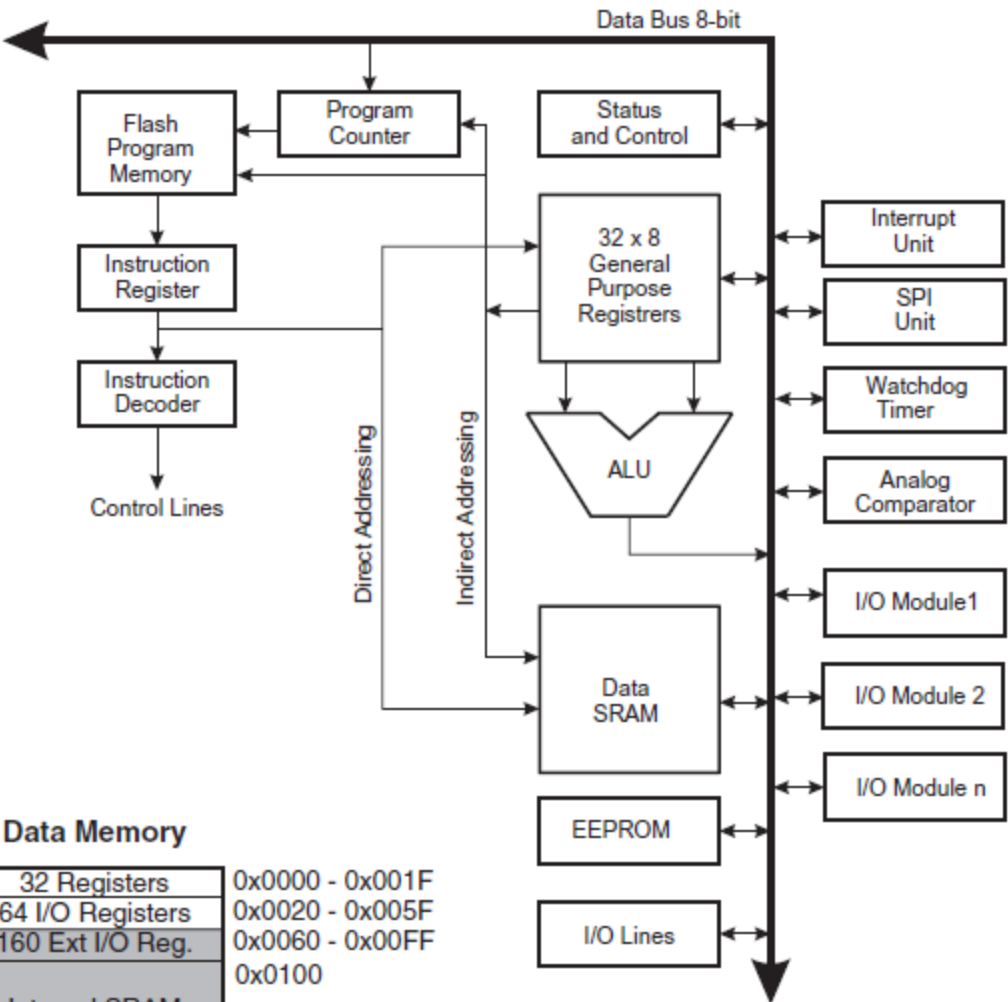
What is meant by the following C code:

```
char *x;
x = 0x20;
```

Data Memory

32 Registers	0x0000 - 0x001F
64 I/O Registers	0x0020 - 0x005F
160 Ext I/O Reg.	0x0060 - 0x00FF
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Understanding the memory architecture



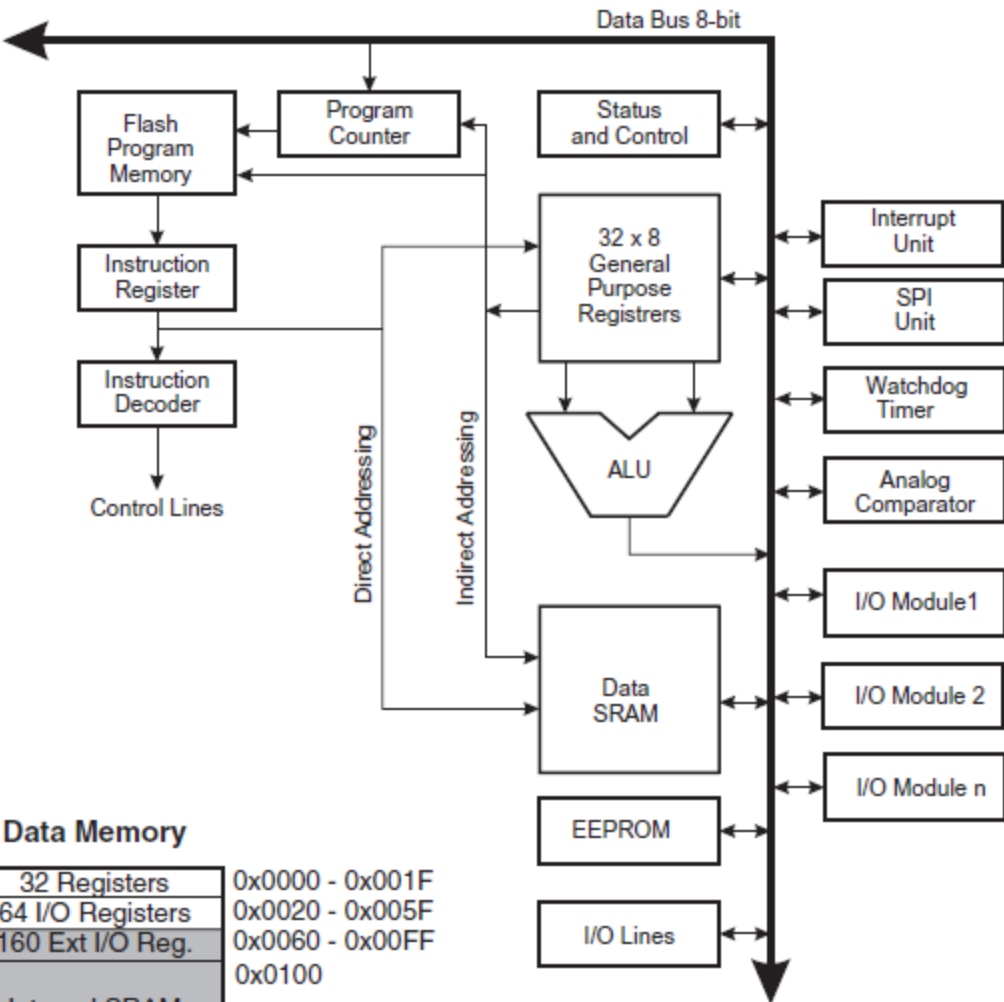
What is meant by the following C code:

```
char *x, y;
x = 0x20;
y = *x;
```

Data Memory

32 Registers	0x0000 - 0x001F
64 I/O Registers	0x0020 - 0x005F
160 Ext I/O Reg.	0x0060 - 0x00FF
Internal SRAM (512/1024/1024 x 8)	0x0100 0x02FF/0x04FF/0x04FF

Understanding the memory architecture



What is meant by the following C code:

```
char *x, y;
x = &y;
*x = 0x20;
```

Data Memory

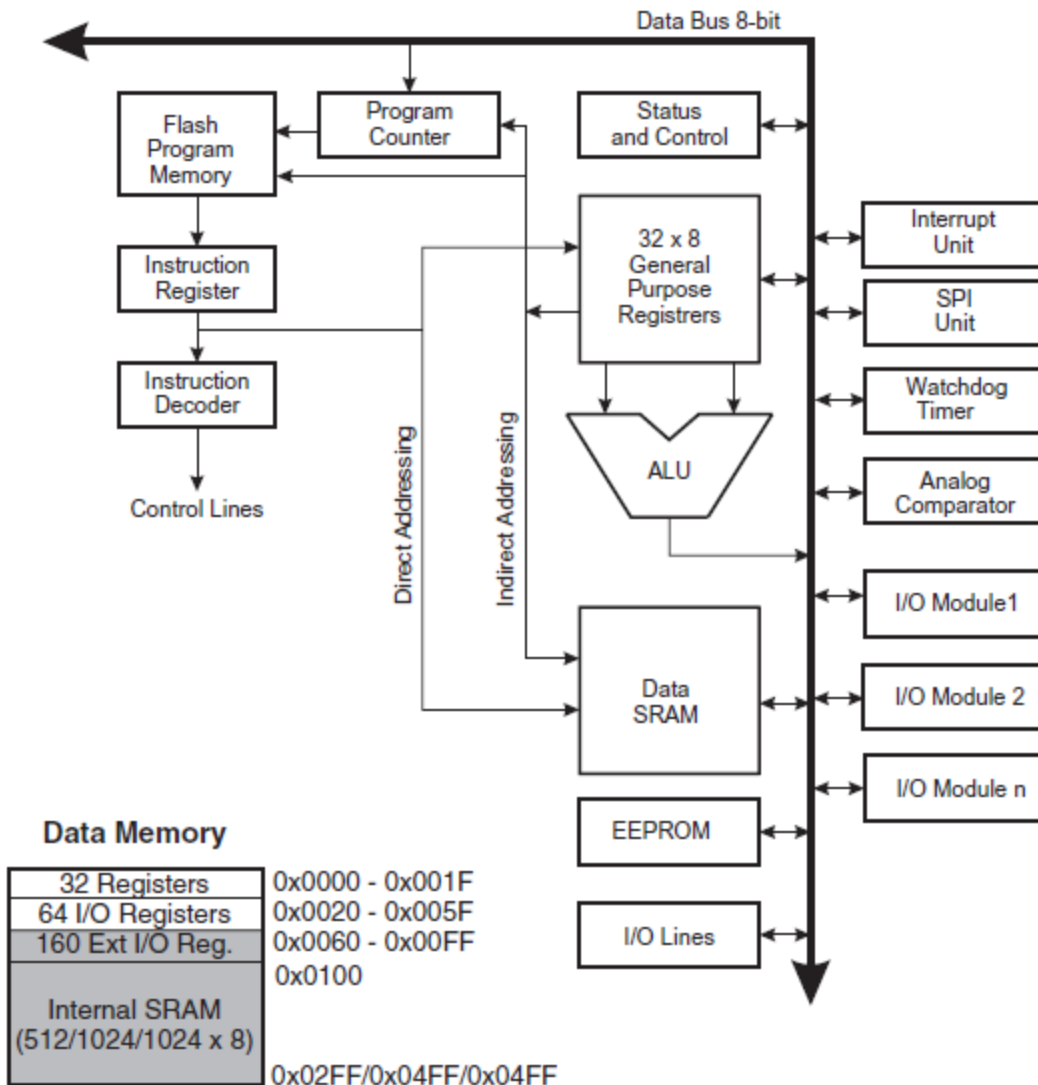
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Understanding the memory architecture

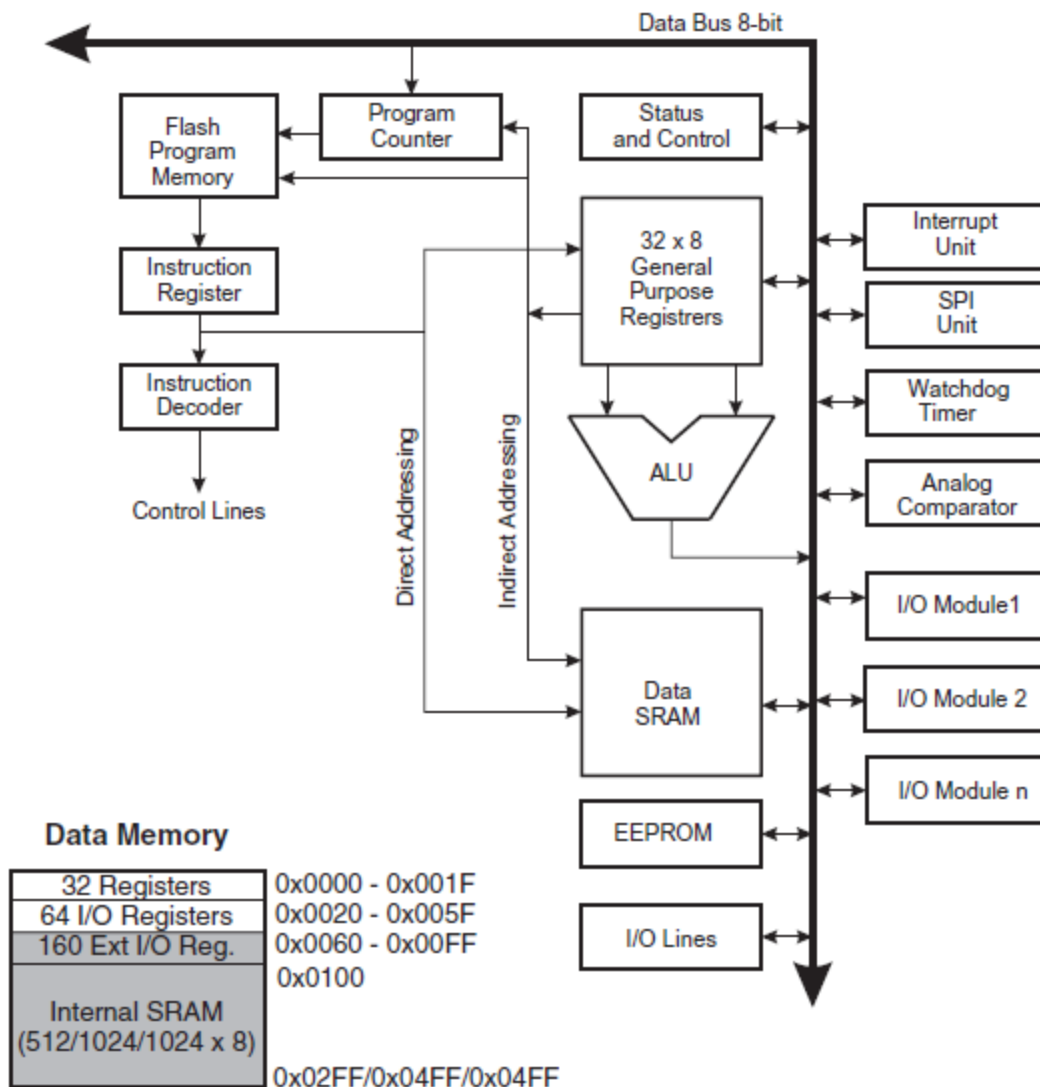
What is meant by the following C code:

```
char foo() {
    char *x, y;
    x = 0x20;
    y = *x;
    return y;
}
char z;
int main(void) {
    z = foo();
    ...
}
```

Where are x, y, z in memory?



Understanding the memory architecture



```
char foo() {
    char *x, y;
    x = 0x20;
    y = *x;
    return y;
}
```

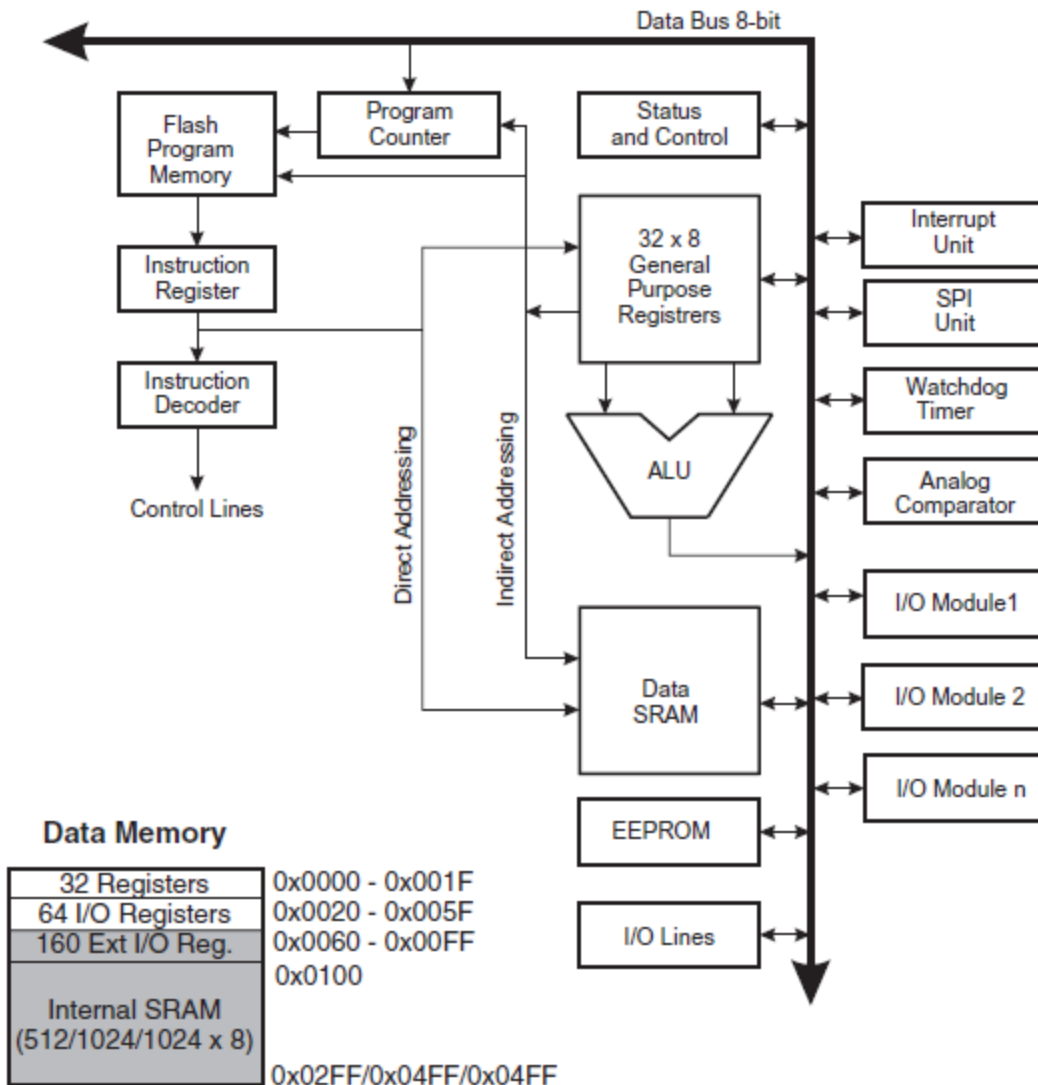
```
char z;
int main(void) {
    z = foo();
    ...
}
```

x occupies 2 bytes on the stack, y occupies 1 byte on the stack, and z occupies 1 byte in static memory.

Understanding the memory architecture

What is meant by the following C code:

```
char foo() {
    char y;
    uint16_t x;
    x = 0x20;
    y = *x;
    return y;
}
char z;
int main(void) {
    z = foo();
    ...
}
```



Memory usage: Understanding the stack.

Find the flaw in this program

```
int x = 2;
```

statically allocated: compiler assigns a memory location.

```
int* foo(int y) {
```

arguments on the stack

```
int z;
```

automatic variables on the stack

```
z = y * x;
```

```
return &z;
```

```
}
```

```
int main(void) {
```

```
int* result = foo(10);
```

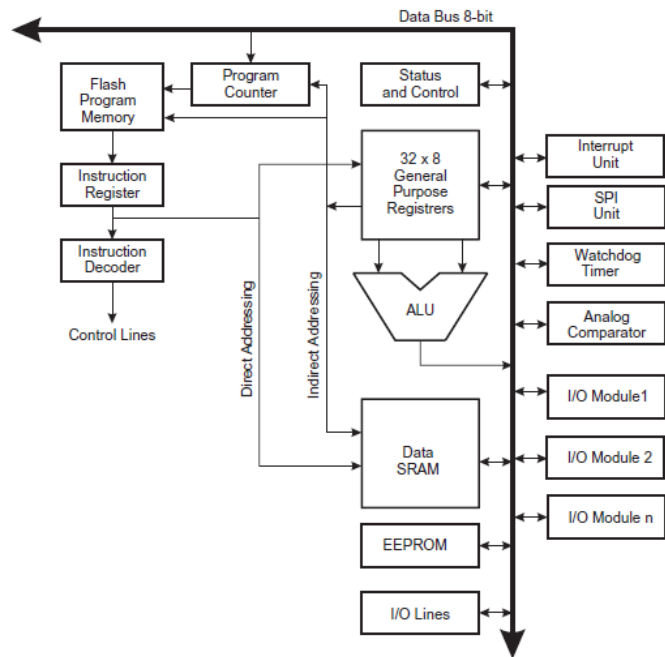
program counter and copies of all registers on the stack

```
...
```

```
}
```

This program returns a pointer to a variable on the stack. What if another procedure call occurs before the returned pointer is de-referenced?

Understanding the memory architecture



```
void foo(uint16_t x) {
    char y;
    y = *x;
    if (x > 0x100) {
        foo(x - 1);
    }
}
char z;
void main(...) {
    z = 0x10;
    foo(0x04FF);
    ...
}
```

What is the value of z?

Data Memory

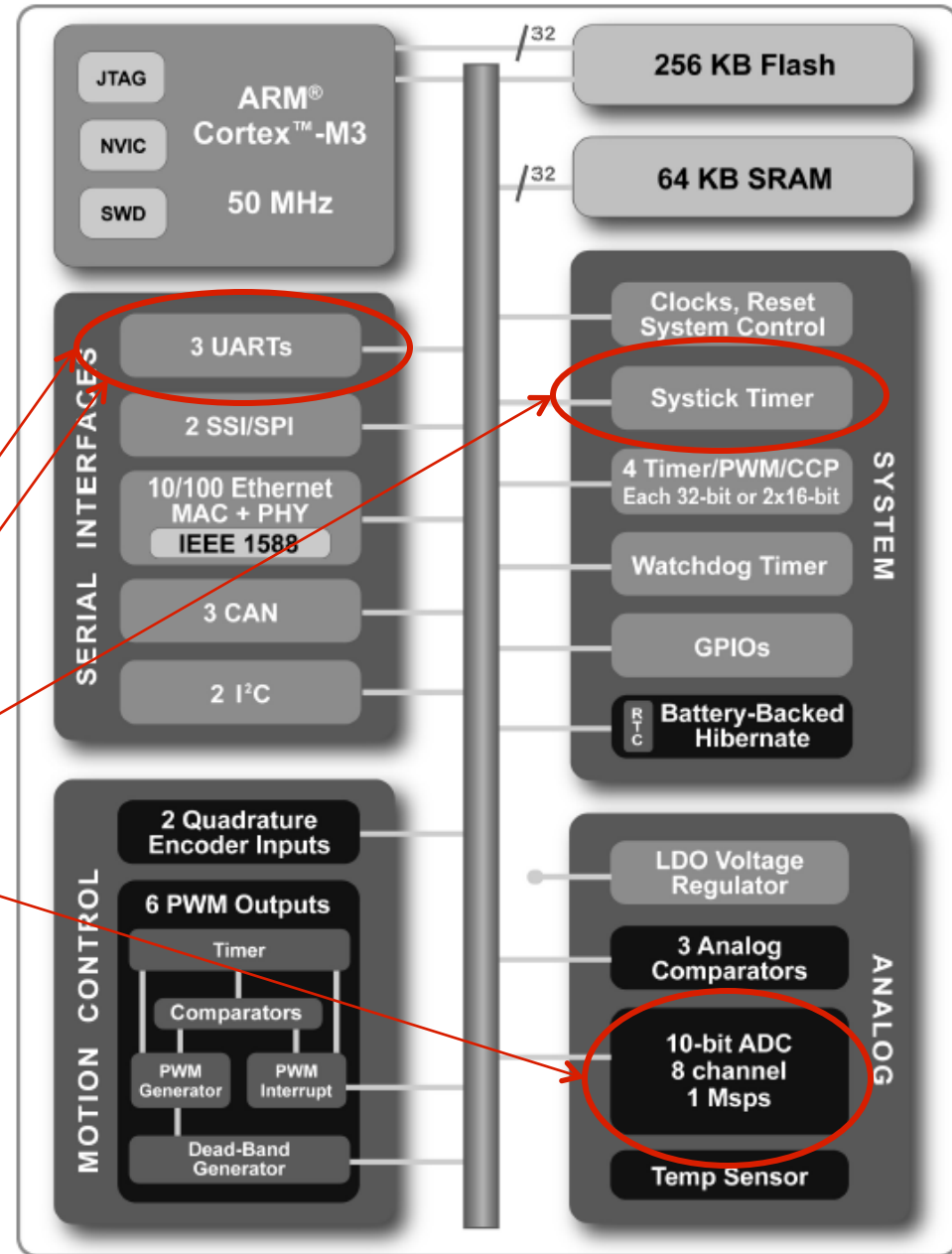
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Luminary LM3S8962 I/O Architecture

32-bit microcomputer.

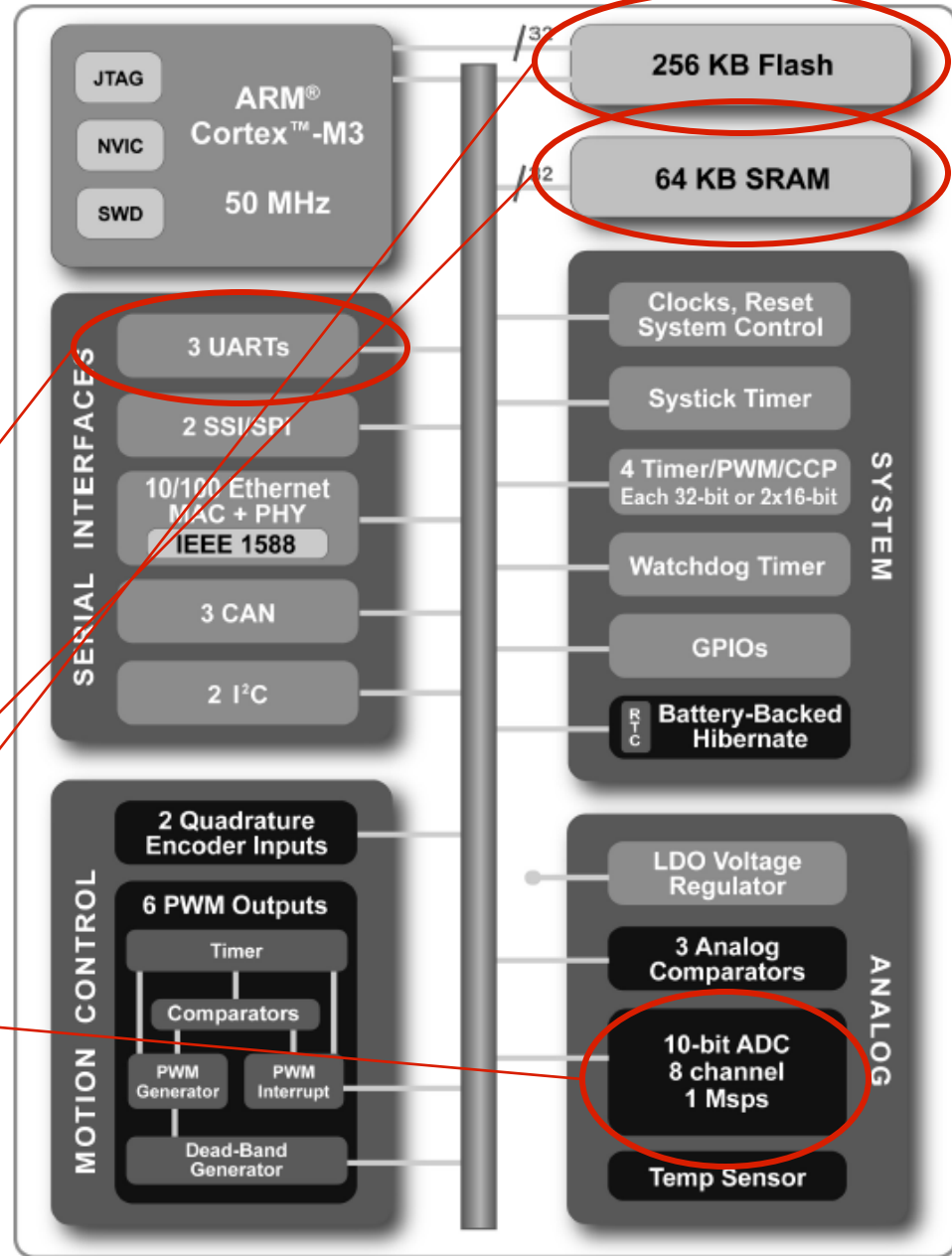
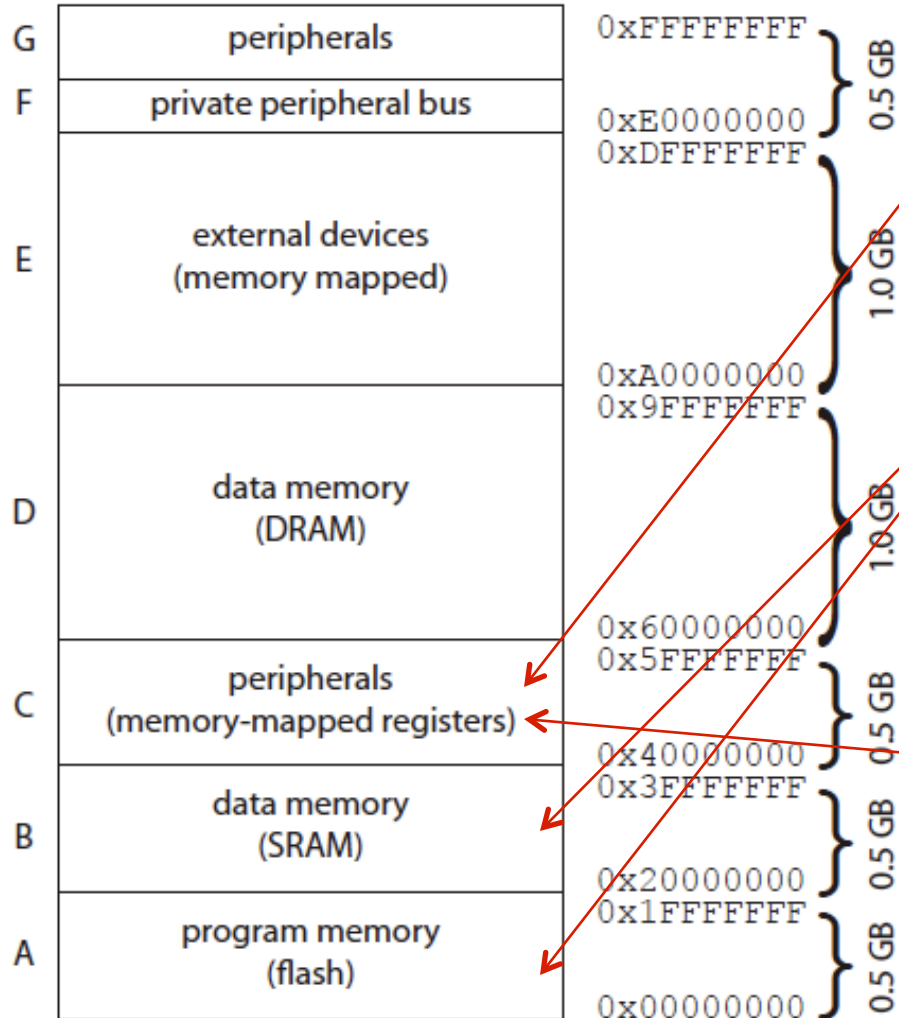
Relevant features:

- UART interface
- ADC to accelerometer
- SysTick/other Timers
- UART to Bluetooth
- JTAG interface for programming
- *No floating point!*



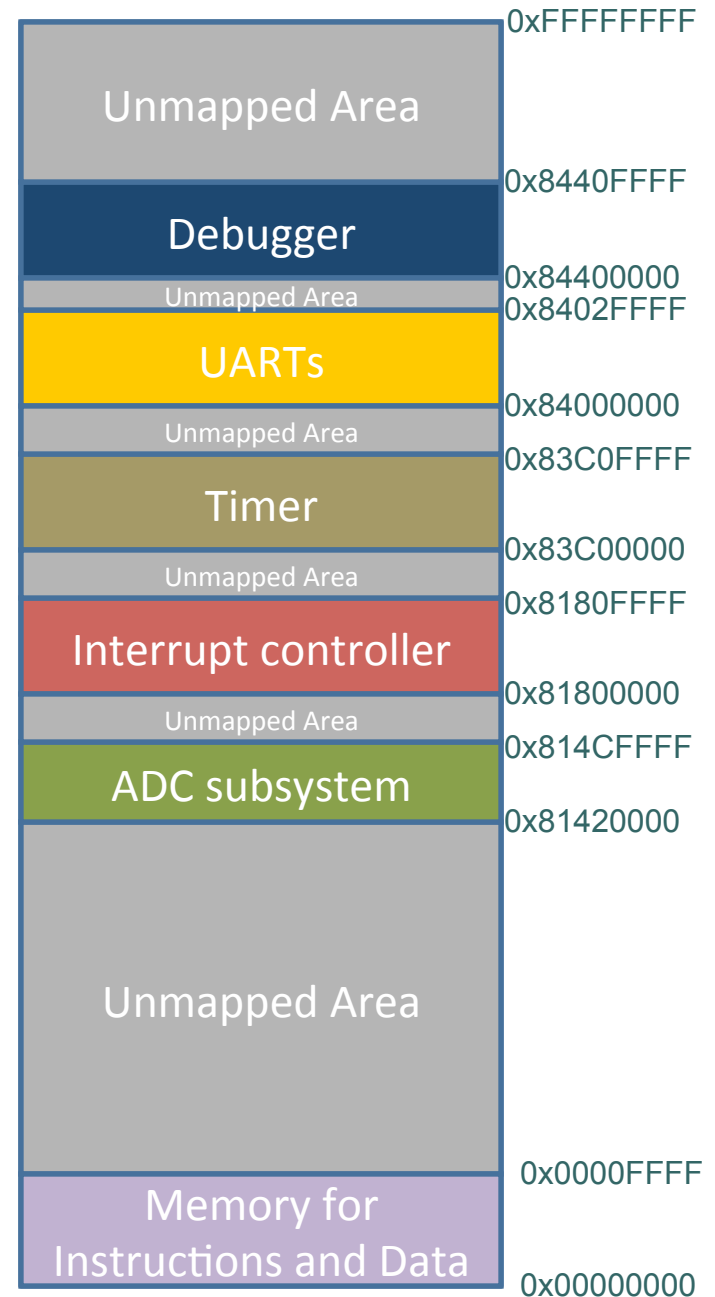
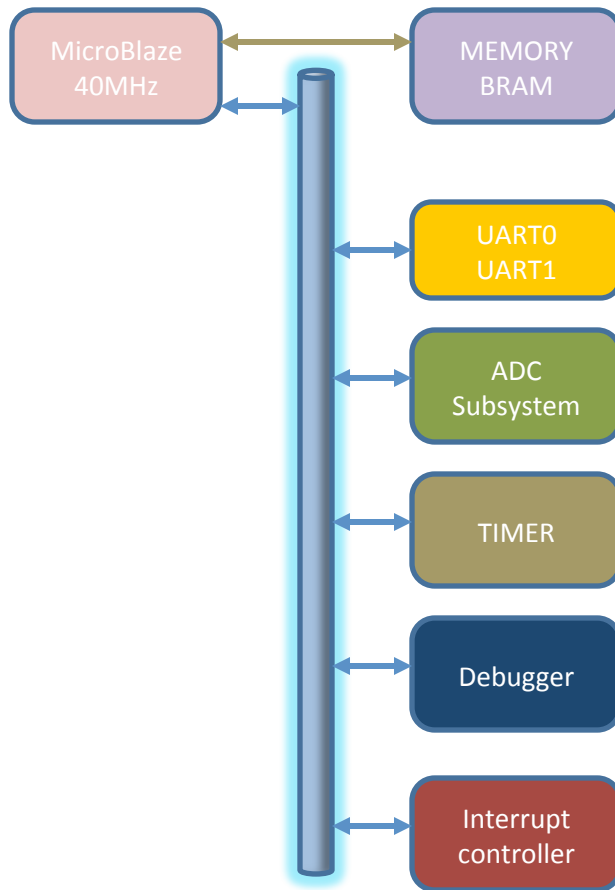
Memory Architecture

ARM Cortex M3



Microblaze memory map

Xilinx soft core



Thanks to Trung Tran

Heaps

An operating system typically offers a way to dynamically allocate memory on a “heap”.

Memory management (`malloc()` and `free()`) can lead to many problems with embedded systems:

- Memory leaks (allocated memory is never freed)
- Memory fragmentation (allocatable pieces get smaller)

Automatic techniques (“garbage collection”) typically require stopping everything and reorganizing the allocated memory. This is deadly for real-time programs.

Summary

Understanding memory architectures is essential to programming embedded systems.